## The Flash code

- framework of many different modules
$\triangleright$ grid-types (AMR, uniform)
$\triangleright$ Rieman solvers (PPM, van Leer, . . . )
$\triangleright$ time-integrators (euler1, rk3, strang)
$\triangleright$ physical modules (self-gravity, MHD, radiation, turbulence stirring, nuclear burning, . . . )
$\triangleright$ particle modules (N-body, SPH)
$\triangleright$ multiple fluids
- modules quite independent, well defined interface
- AMR - block-type, based on PARAMESH library
$\triangleright$ relatively easy to code
$\triangleright$ can completely automatic (different refinement criteria)
$\triangleright$ less efficient
- self-gravity in Flash
$\triangleright$ multipole solver (bad scaling with \#CPU, must be close to spherical symmetry)
$\triangleright$ multigrid solver (very slow - lot of communication, artifacts at different AMR levels)
$\triangleright$ Why not tree code?


## AMR in Flash



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## Load balancing - Morton curve



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## Load balancing - Morton curve



Gravity tree


## Gravity tree



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Gravity tree


## The code scheme

gr_solversInit - initialization, allocation of arrays, setting constants, etc.

Grid_solvePoisson (in each time-step):
gr_treeBuildTree $\longleftarrow<$ gr_treeBuildTreeBlock $\Longleftarrow\left\{\begin{array}{l}\text { gr_treeGetTreePos } \\ \text { gr_treeGetTreeSize }\end{array}\right.$ gr_treeExchangeTrees
gr_treePotential $<$ gr_treePotentialBlock
gr_treeDestroyTree
gr_solversFinalize - deallocation of arrays

## Interfaces:

gr_treeData - global variables of the module gr_treeInterface - prototypes of subroutines

## Tree in RAM


tree nodes identified by multi-index - integer array of size $L:\left(l_{1}, l_{2}, l_{3}\right) ; l_{i}=$
$\triangleright 1-8 \ldots$ number of node on $i$-th level
$\triangleright 0 \ldots$ multi-index (i.e. node) is of level i-1

## Multi-index - ID of the node


coordinates $\rightarrow$ node multi-index

```
do l = 1, tr_levels
    fac = 2**(tr_levels - l)
    multi(l) = 1 + mod((i-1), 2*fac)/fac &
    +2*(mod}((j-1),2*fac)/fac) &
    + 4 *(mod((k-1),2*fac)/fac)
```

enddo
node multi-index $\rightarrow$ coordinates
$(i, j, k)=(1,1,1)$
do $1=1$,tr_levels
fac $=2 * *($ tr_levels - 1 )
$i=i+(\bmod (\operatorname{multi}(1)-1,2)) \quad *$ fac
$j=j+(\bmod ((\operatorname{multi}(1)-1) / 2,2)) *$ fac
$\mathrm{k}=\mathrm{k}+(\bmod ((\operatorname{multi}(1)-1) / 4,2))$ * fac enddo

## Linear index in the tree array

```
integer function get_tree_pos(level, mi)
    use gr_treeData, ONLY : tr_levels
    integer,intent(in) :: mi(1:tr_levels)
    integer :: pos, fs
    ! field size: 1 for the lowest level (only masses),
    ! 4 for the higher levels (mass, position of mass centre)
    if (level == tr_levels) then
        fs = 1
    else
        fs = 4
    endif
    pos = 1 + 4*(8**level - 1)/7 ! offset of the level
    do l = level,1,-1
        pos = pos + (mi(l)-1)*fs
        fs = fs * 8
    enddo
    get_tree_pos = pos
end function get_tree_pos
```


## Building tree

- create the lowest level (3)
$\triangleright$ last segment of tree array filled with cell masses
$\triangleright$ the 2nd lowest level prepared $-m_{i}, m_{i} \times \mathbf{r}_{i}$ accumulated
- finish the 2nd lowest level
$\triangleright \mathbf{r}_{m c}=\sum m_{i} \times \mathbf{r}_{i} / m$
- create the higher levels
$\triangleright$ loop over levels (second lowest and higher)
$\triangleright$ loop over nodes of the given level
$\triangleright$ collect contribution from each octet and write it into parent node


## Communication - data structures

- communication arrays that are synchronized among all CPUs integer tr_activeBlocks(MAXBLOCKS, \#CPUs) - 1 denotes LEAF block real tr_BBoxes(MAXBLOCKS, 2, DIM, \#CPUs) - bounding boxes real tr_Coords(MAXBLOCKS, 2, BLOCK_SIZE, \#CPUs) - coordinates real tr_Diag(tr_levels, MAXBLOCKS, \#CPUs) - node sizes real tr_sentTreeLevels(MAXBLOCKS, \#CPUs (to), \#CPUs (from))
- trees - array of pointers
! array of pointers to trees $2 \mathrm{D}=$ (\#CPUs, MAXBLOCKS)
type p_tree
real, dimension(:), pointer :: p
end type p_tree
type(p_tree), save, dimension(:,:), allocatable :: tree_array
- communication costs:
$\triangleright$ ping roundtrip on coma $\sim 0.2 \mathrm{~ms}$
( $\rightarrow$ establishing TCP connection $\sim 1 \mathrm{~ms}$ )
$\triangleright 1 G b$ ethernet $=128 \mathrm{MB} / \mathrm{s} \Rightarrow 128 \mathrm{~KB} / \mathrm{ms}$
$\triangleright$ establishing connection $\Leftrightarrow$ transmitting array with 16000 elements


## Communication scheme

- synchronize global block describing arrays
tr_activeBlocks, tr_BBoxes, tr_Coords, tr_Diag
- determine tr_sentTreeLevels
- compare BBox of each block with BBoxes of all blocks on to-CPU
$\triangleright$ searches for the minimum distance of corners (cmp. all pairs)
- synchronize tr_sentTreeLevels
- prepare messages
$\triangleright$ for each CPU 1D array which contains trees to the necessary level
$\triangleright$ allocate arrays for messages to be recieved
- exchange messages - non-blocking communication (MPI_Isend, MPI_recv)
- copy trees from messages to tree_array


## Potential computation

- for each point



## Accuracy test



## Performance



## Scaling with \# of CPU



## Scaling with \# of blocks

time on 2 CPUs, 100 time-steps, $\Theta_{\text {lim }}=1.0$


## What next

- higher tree levels (to be less bush-like :))
$\triangleright$ in theory simple: just to add $m, \mathbf{r}_{m c}$ for non-LEAF blocks
- some optimalization ( $1 / x$, inline function)
- quadrupole moments
- any ideas?

